Towards a Theory of Anaphoric Binding in Event Semantics

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Outline

▶ Introduction

Polynomial Event Semantics

Relative Denotations

Nominal Pronouns and Referents

Conclusions

Give compositional denotations and Decide entailments (in FraCaS, etc.) with Neo-Davidsonian event semantics

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- ▶ Quantification: some, all, every
- ▶ Complex quantification: many, most, few, at least 5, ...
- Negative quantification and Negation
- Copula clauses
- Relative clauses
- ▶ Anaphora: pronouns, crossover, ellipsis, ...

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- 6 Complex quantification: many, most, few, at least $5, \ldots$
- 4 Negative quantification and Negation
- 3 Copula clauses
- 2 Relative clauses
- 1 Anaphora: pronouns, crossover, ellipsis, ...

Rank in importance

Give compositional denotations and Decide entailments (in FraCaS, etc.) with Neo-Davidsonian event semantics

- 1 Quantification: a, some, all, every
- 3 Complex quantification: many, most, few, at least $5, \ldots$
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- 4 Copula clauses
- 5 Relative clauses
- 6 Anaphora: pronouns, crossover, ellipsis, ...

Rank in attention received in event semantics research

Polynomial Event Semantics

A *variable-free* dialect of Neo-Davidsonian event semantics specifically designed to address the problems

- Quantification: some, all, every LENLS 2018
- Complex quantification: many, most, few, at least 5, ... LENLS 2021
- Negative quantification and Negation LENLS 2020
- Copula clauses LENLS 2021
- Relative clauses LENLS 2022
- Anaphora: pronouns, crossover, ellipsis, ...
 Beginning now

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Anaphoric binding becomes oddly symmetric

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 $[[John traveled to Paris.]] = (subj' / john) \cap Travel \cap TP$

 $[[John traveled to Paris.]] = (subj' / john) \cap Travel \cap TP$

The meaning of the whole sentence is the intersection of the meaning of its constituents

[John traveled to Paris.]]= (subj'/ john) \cap Travel \cap TP

The denotation of a sentence

- ▶ a set of events that witness it
- ▶ a formula that represents the set
- ▶ a query of the record of world events

The sentence is true in a world if the result of the query is non-empty

[Bill and John traveled to Paris.]] = $subj'/(john \otimes bill) \sqcap Travel \sqcap TP$

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Poly-Denotations: Further Choices

$$[Bill or John]] = \begin{cases} bill \oplus john & \text{if external} \\ bill \sqcup john & \text{if internal} \end{cases}$$

Poly-Denotations: Quantification

$$\llbracket \text{Everyone} \rrbracket = \bigotimes_{i \in \text{Person}} i = \mathcal{A} \text{Person}$$
$$\llbracket A_W \text{ person} \rrbracket = \bigoplus_{i \in \text{Person}} i = \mathcal{I} \text{ Person}$$
wide-scope; indefinite

$$\llbracket A_N \text{ person} \rrbracket = \bigsqcup_{i \in \mathsf{Person}} i = \mathcal{E} \operatorname{Person}_{\operatorname{narrow-scope; some}}$$

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the denotation was compositionally derived

$$[T] =?$$

Generalization to other anaphora

Relative *things*

Things d

individuals, sets (concepts), poly-concepts, already relitivized individuals and concepts, etc.

Relative things

relations between individuals and things, or sets of pairs (i, d)

$$\{(i, d) \mid i \in C\}$$

where C is some set
$$\equiv d|i:C$$

Relativitizing

$$d \stackrel{\iota}{\underset{\rho}{\rightleftarrows}} d|i:C$$

where

$$\begin{split} \iota_C \; d &= d | i : C \\ & \text{inclusion (or, embedding)} \\ \rho \; (d | i : C) &= d \\ & \text{provided } d \text{ is independent of } i \\ & \text{retract (or, projection)} \end{split}$$

Relativitizing

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where

$$\iota_C d = d|i:C$$

inclusion (or, embedding)
 $\rho (d|i:C) = d$ provided d is independent of i
retract (or, projection)

 ι and ρ are inverses, but not a bijection: not all relations are constant

Relativitizing

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When C is singleton, ι and ρ are a bijection: A singleton context may *always* be disposed of

Denotations so far
individualj
concept (event set)d
constant \bot Operations $\oplus, \otimes, \sqcap, \sqcup$ Algebra of poly-concepts (poly-individuals)

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$$(d_1|i:C) \sqcap (d_2|i:C) = (d_1 \sqcap d_2)|i:C$$

Relative denotations

individual	j i:C
concept (event set)	d i:C
constant	$\perp i:C$
Operations	$\oplus, \otimes, \sqcap, l$

Denotations so far
individualjconcept (event set)dconstant \bot Operations $\oplus, \otimes, \sqcap, \sqcup$ Algebra of poly-concepts (poly-individuals)

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Relative denotations
individual j|i:C
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Operations $\oplus, \otimes, \sqcap, \sqcup$ Algebra of relative poly-concepts (poly-individuals)

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the denotation was compositionally derived

Traces

$$\begin{bmatrix} \text{city John traveled to } T \end{bmatrix} \\ = \text{City} \cap \{i \mid \llbracket \text{John traveled to } i \rrbracket \neq \bot \}$$

the denotation was *compositionally derived*

$$\llbracket T \rrbracket = i | i: \mathcal{I}$$

(\mathcal{I} : set of all individuals)

$[\![It is famous.]\!]$

[It is famous.]

[[it]] = ?

[It is famous.]

 $\llbracket it \rrbracket = i | i: Thing$

$[[It is famous.]] = (subj' / i \cap Famous) | i: Thing$

(Embeddings are applied as necessary, silently)

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 $[\![It is famous.]\!] \sim [\![Who is famous?]\!]$

$[\![John \ traveled \ to \ Paris^{\triangleright}.]\!]$

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$[\![\operatorname{Paris}^{\triangleright}]\!]$

$[\![John \ traveled \ to \ Paris^{\triangleright}.]\!]$

$$[[Paris^{\triangleright}]] = \iota_{\{paris\}} \text{ paris}$$

[John traveled to Paris[▷].]]

$[Paris^{\triangleright}]]$ = $\iota_{\{paris\}}$ paris = paris |*i*: {paris}

$[\![John \ traveled \ to \ Paris^{\triangleright}.]\!]$

$[Paris^{\triangleright}]]$ $= \iota_{\{paris\}} paris$ $= paris |i: \{paris\}$ $= i |i: \{paris\}$

$[\![John \ traveled \ to \ Paris^{\triangleright}.]\!]$

$$[\operatorname{Paris}^{\triangleright}] = i | i: \{ \mathsf{paris} \}$$
$$[it] = i | i: \mathsf{Thing}$$

$$\begin{split} & [\text{John traveled to Paris}^{\triangleright}.] \\ &= [\text{John traveled to } i] | i: \{ \text{paris} \} . \end{split}$$

$$[\operatorname{Paris}^{\triangleright}] = i | i: \{ \text{paris} \}$$
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The denotation is not inherently relative (ρ does apply): contains no unresolved anaphoric references

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$$[\operatorname{Paris}^{\triangleright}] = i | i: \{ \text{paris} \}$$
$$[it] = i | i: \text{Thing}$$

- The denotation is not inherently relative (ρ does apply): contains no unresolved anaphoric references
- ▶ Referent and reference: odd symmetry

Discourse and Resolution

John traveled to Paris^{\triangleright}. It is famous.

$$(d_1|i:C) \sqcap (d_2|i:C) = (d_1 \sqcap d_2)|i:C$$

 $(d_1|i:C_1) \sqcap (d_2|i:C_2) = ? \qquad (C_1 \neq C_2)$

$\iota_{C_2}(d_1|i_1:C_1) \ \sqcap \ \iota_{C_1}(d_2|i_2:C_2)$

$$\begin{split} \iota_{C_2}(d_1|i_1:C_1) &\sqcap \ \iota_{C_1}(d_2|i_2:C_2) \\ &= (d_1|i_1:C_1)|i_2:C_2 \ \sqcap \ (d_2|i_2:C_2)|i_1:C_1 \end{split}$$

$$\begin{aligned}
\iota_{C_2}(d_1|i_1:C_1) &\sqcap \iota_{C_1}(d_2|i_2:C_2) \\
&= (d_1|i_1:C_1)|i_2:C_2 &\sqcap (d_2|i_2:C_2)|i_1:C_1 \\
&= (d_1|(i_1:C_1 \times i_2:C_2)) \sqcap (d_2|(i_1:C_1 \times i_2:C_2))
\end{aligned}$$

$$\begin{split} \iota_{C_2}(d_1|i_1:C_1) &\sqcap \iota_{C_1}(d_2|i_2:C_2) \\ &= (d_1|i_1:C_1)|i_2:C_2 &\sqcap (d_2|i_2:C_2)|i_1:C_1 \\ &= (d_1|(i_1:C_1 \times i_2:C_2)) \sqcap (d_2|(i_1:C_1 \times i_2:C_2)) \\ &= (d_1 \sqcap d_2)|(i_1:C_1 \times i_2:C_2) \end{split}$$

Discourse and Resolution

 $\llbracket John traveled to Paris^{\triangleright}.$ It is famous. \rrbracket

= (\llbracket John traveled to $i_1 \rrbracket \otimes \llbracket i_2$ is famous \rrbracket)| $(i_1: \{paris\} \times i_2: Thing)$

 $\llbracket John traveled to Paris^{\triangleright}.$ It is famous. \rrbracket

= (\llbracket John traveled to $i_1 \rrbracket \otimes \llbracket i_2$ is famous \rrbracket) $|(i_1: \{paris\} \times i_2: Thing)$

{Imposing an equality constraint via pragmatics.} \Rightarrow ([John traveled to i] \otimes [i is famous])|i : {paris} $\llbracket John \ traveled \ to \ Paris^{\triangleright}.$ It is famous. \rrbracket

= (\llbracket John traveled to $i_1 \rrbracket \otimes \llbracket i_2$ is famous \rrbracket) $|(i_1: \{paris\} \times i_2: Thing)$

{Imposing an equality constraint via pragmatics.}

- $\Rightarrow (\llbracket \text{John traveled to } i \rrbracket \otimes \llbracket i \text{ is famous} \rrbracket) | i : \{ \mathsf{paris} \}$
- $\stackrel{\rho}{\rightarrow} \llbracket \text{John traveled to Paris.} \rrbracket \otimes \llbracket \text{Paris is famous.} \rrbracket$

The result after resolution is 'context-free'

Time Arrow

Are things too symmetric? John travel to it. Paris is famous.

Anaphora vs. Cataphora

$$\begin{aligned} \iota_{C_2}(d_1|i_1:C_1) &\sqcap \iota_{C_1}(d_2|i_2:C_2) \\ &= (d_1|i_1:C_1)|i_2:C_2 &\sqcap (d_2|i_2:C_2)|i_1:C_1 \\ &= (d_1|(i_1:C_1 \times i_2:C_2)) \sqcap (d_2|(i_1:C_1 \times i_2:C_2)) \\ &= (d_1 \sqcap d_2)|(i_1:C_1 \times i_2:C_2) \end{aligned}$$

Context as an Imagination Restriction

Discourse/Context is a constraint on listeners' imagination

- ▶ The more narrow is the context, the tighter is the constraint
- In the limit, a proper noun is the anaphoric reference in the singleton context, which constrains it unambiguously

The close similarity of $[\![\mathrm{Paris}^{\triangleright}]\!]$ and $[\![\mathrm{it}]\!]$ should not be so surprising

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Anaphoric dependencies can, after all, be expressed in a variable-free event semantics in a surprisingly symmetric way: Both the referent and the reference denoted by a polyconcept relative to a context

Context as an Imagination Restriction

The development of the theory of anaphoric binding in event semantics has just began Many, many more examples to analyze